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Semimodule enrichment

Chris Heunen

*Institute for Computing and Information Sciences
Radboud University, Nijmegen, the Netherlands,
Email: heunen@cs.ru.nl*

Abstract

A category with biproducts is enriched over (commutative) additive monoids. A category with tensor products is enriched over scalar multiplication actions. A symmetric monoidal category with biproducts is enriched over semimodules. We show that these extensions of enrichment (*e.g.* from hom-sets to hom-semimodules) are functorial, and use them to make precise the intuition that “compact objects are finite-dimensional” in standard cases.

Keywords: Semimodules, enriched categories, biproducts, scalar multiplication, compact objects.

1 Introduction

Mitchell’s celebrated theorem states that every Abelian category can be embedded in the category of modules over a ring [18]. This article is a first part of a generalisation from rings to semirings¹, which could hopefully give a representation theorem for semantic models of linear (quantum) computation. An important step in the embedding theorem is that a category with biproducts is enriched over commutative monoids. Mitchell’s proof then continues by finding an appropriate scalar multiplication. More recently, Abramsky observed that a category with tensor products always has a natural scalar multiplication [1]. We will prove that tensor products in fact provide enrichment with scalar multiplication. However, Abramsky’s and Mitchell’s scalars differ. We combine both to show that a category that has biproducts as well as tensor products is enriched over semimodules [7].² In fact, we show that this extension of enrichment (from **Set** to semimodules) is functorial, and holds for any enriching category **V** instead of **Set**. By way of introduction, let us discuss these results in the case of ordinary categories briefly.

Biproducts give additive enrichment

Recall that a *zero object* is an object that is simultaneously initial and terminal. Its existence means that there is a unique morphism $0_{XY} : X \rightarrow Y$ factoring through

¹ A semiring is roughly a ring that does not necessarily have subtraction.

² A semimodule is to a semiring what a module is to a ring, and a vector space to a field.

the zero object between any two objects X and Y . A category \mathbf{C} is said to have *binary biproducts* when it has coinciding (binary) products and (binary) coproducts such that

$$\nabla_{X_1 \oplus X_2} \circ ((\kappa_1 \circ \pi_1) \oplus (\kappa_2 \circ \pi_2)) \circ \Delta_{X_1 \oplus X_2} = \text{id}_{X_1 \oplus X_2}, \quad (1)$$

$$\pi_i \circ \kappa_i = \text{id}_{X_i}, \quad (2)$$

$$\pi_j \circ \kappa_i = 0_{X_i X_j} \text{ when } i \neq j, \quad (3)$$

where we write $X_1 \oplus X_2$ for $X_1 \times X_2 = X_1 + X_2$, Δ for the diagonal (id, id) , ∇ for the codiagonal $[\text{id}, \text{id}]$, and π and κ for the projections and coprojections. A category has *finite biproducts* when it has a zero object and binary biproducts. In order to prepare for later generalisation we shall be a bit formal and write **Set-Cat** for the category of **Set**-enriched categories, *i.e.* of locally small categories. We denote by **BP(Set-Cat)** the category of all locally small categories with finite biproducts and functors preserving them. By **cMon(C)** we denote the category of all commutative monoids in a symmetric monoidal category \mathbf{C} ; in case $\mathbf{C} = \mathbf{Set}$, we abbreviate it to **cMon**. Section 2 considers an extension of enrichment that for ordinary categories is given as follows.

Theorem 1.1 *Locally small categories with finite biproducts are cMon-enriched, and this is functorially so: there is a functor $\mathbf{BP}(\mathbf{Set-Cat}) \rightarrow \mathbf{cMon-Cat}$.*

Proof. We describe the monoid structure on the homset $\mathbf{C}(X, Y)$ additively. The sum $f + g : X \rightarrow Y$ of $f : X \rightarrow Y$ and $g : X \rightarrow Y$ is given by

$$f + g : X \xrightarrow{\Delta} X \oplus X \xrightarrow{f \oplus g} Y \oplus Y \xrightarrow{\nabla} Y. \quad (4)$$

The monoid unit is the zero morphism $0_{XY} : X \rightarrow Y$. Since morphisms in **BP(Set-Cat)** preserve biproducts, they also preserve this enrichment. \square

Tensor products give scalar multiplication enrichment

In any monoidal category \mathbf{C} , the endomorphisms of the monoidal unit can be considered as a monoid of scalars, since one can define an action of it on the category called *scalar multiplication* as follows [1]. For a *scalar* $s : I \rightarrow I$ and any morphism $f : X \rightarrow Y$ of \mathbf{C} , define $s \bullet f : X \rightarrow Y$ as

$$X \xrightarrow{\cong} I \otimes X \xrightarrow{s \otimes f} I \otimes Y \xrightarrow{\cong} Y.$$

In particular, $\mathbf{C}(I, I)$ forms a monoid with id_I as unit, and $s \bullet t$ as multiplication. Thus we have the following lemma. For a fixed monoid M in \mathbf{C} , we denote by $\mathbf{Act}_M(\mathbf{C})$ the category of (left) actions of M [15, Section VII.4]; we use the shorthand $\mathbf{Act}_M = \mathbf{Act}_M(\mathbf{Set})$.

Lemma 1.2 *If \mathbf{C} is a locally small monoidal category, then $\mathbf{C}(I, I)$ acts upon its homsets. In other words, $\mathbf{C}(X, Y) \in \mathbf{Act}_{\mathbf{C}(I, I)}$.*

Proof. To verify that the scalar multiplication defined above is an action one has to check that $\text{id}_I \bullet f = f$ and $s \bullet (t \bullet f) = (s \bullet t) \bullet f$. \square

This is the first step towards proving that every symmetric monoidal category \mathbf{C} is enriched over $\mathbf{Act}_{\mathbf{C}(I,J)}$. To complete the reasoning we need to ensure that composition is a morphism in $\mathbf{Act}_{\mathbf{C}(I,J)}$. Because $s \bullet (g \circ f) = g \circ (s \bullet f) = (s \bullet g) \circ f$, we first have to equip $\mathbf{Act}_{\mathbf{C}(I,J)}$ with a tensor structure, universal for bimorphisms. For this it suffices that \mathbf{C} is symmetric monoidal. This will be proven in detail in Section 3 using techniques by Kock and Day. For now, let us just state the result without proof. We write $\mathbf{Act} = \mathbf{Act}(\mathbf{Set})$ for the category of (left) actions of an arbitrary monoid, to be explained in more detail later.

Theorem 1.3 *A symmetric monoidal \mathbf{Set} -category \mathbf{C} is enriched over $\mathbf{Act}_{\mathbf{C}(I,J)}$. This is functorially so: there is a functor $\mathbf{cMon}(\mathbf{Set-Cat}) \rightarrow \mathbf{Act-Cat}$.*

Semimodule enrichment

If \mathbf{C} is a symmetric monoidal category, then so is $\mathbf{cMon}(\mathbf{C})$. Hence one can consider monoid objects in the latter category. First of all, such a ‘double monoid’ object is an object of $\mathbf{cMon}(\mathbf{C})$, *i.e.* a commutative monoid that we write additively as $(S, +, 0)$. Furthermore, it carries monoid structure $I \xrightarrow{1} S \xleftarrow{\cdot} S \otimes S$. Because the latter are morphisms in $\mathbf{cMon}(\mathbf{C})$, we can recognize the objects of $\mathbf{Mon}(\mathbf{cMon}(\mathbf{C}))$ as $S \in \mathbf{C}$ equipped with a commutative monoid structure $(+, 0)$ and a monoid structure $(\cdot, 1)$ such that

$$\begin{aligned} s \cdot (s' + s'') &= (s \cdot s') + (s \cdot s''), \\ (s + s') \cdot s'' &= (s \cdot s'') + (s' \cdot s''), \\ s \cdot 0 &= 0 = 0 \cdot s. \end{aligned}$$

In other words: $\mathbf{Mon}(\mathbf{cMon}) = \mathbf{SRng}$, the category of semirings. This identification allows one to consider a *semiring* whose carrier is not a set: we define $\mathbf{SRng}(\mathbf{C}) = \mathbf{Mon}(\mathbf{cMon}(\mathbf{C}))$ for any symmetric monoidal category \mathbf{C} . Although the multiplication (*i.e.* the outer monoid multiplication \cdot) need not be commutative in general, we will mostly consider only the full subcategory $\mathbf{cSRng}(\mathbf{C})$ of commutative semirings.

Given a monoid object S in a symmetric monoidal category $\mathbf{cMon}(\mathbf{C})$, one can consider the actions of S on objects of the category. In a similar fashion as above, we can recognize these as semimodules over the semiring S ; the commutative monoid that is being acted upon then is the additive group of the semimodule. After all, a semimodule is a space of ‘vectors’ that can be added, and can be multiplied by a scalar [7]. In other words: $\mathbf{Act}_S(\mathbf{cMon}) = \mathbf{SMod}_S$, the category of semimodules over S . This identification allows one to consider a *semimodule* without an underlying set: we define $\mathbf{SMod}_S(\mathbf{C}) = \mathbf{Act}_S(\mathbf{cMon}(\mathbf{C}))$ for any symmetric monoidal category \mathbf{C} .

Let us consider a few examples. It is well-known that \mathbb{Z} -modules are simply Abelian groups, so that $\mathbf{Act}_{\mathbb{Z}}(\mathbf{cMon}) \cong \mathbf{Ab}$. Since an action of \mathbb{N} on a commutative monoid is already completely fixed by its monoid structure, we have $\mathbf{Act}_{\mathbb{N}}(\mathbf{cMon}) \cong \mathbf{cMon}(\mathbf{Set})$. Looking at the Boolean semiring $\mathbb{B} = (\{0, 1\}, \max, 0, \min, 1)$, we can identify an action of \mathbb{B} on a commutative monoid L as an idempotent commutative monoid, because $l = 1 \cdot l = \max(1, 1) \cdot l = \max(1 \cdot l, 1 \cdot l) = \max(l, l)$. Hence $\mathbf{Act}_{\mathbb{B}}(\mathbf{cMon}) \cong \mathbf{SLat}$, the category of (bounded) semilattices.

Now suppose we have a locally small category \mathbf{C} with finite biproducts as well as tensor products, *i.e.* $\mathbf{C} \in \mathbf{cMon}(\mathbf{BP}(\mathbf{Set-Cat}))$.³ Then it is enriched with both addition and scalar multiplication distributing over it. In other words, combining Theorems 1.1 and 1.3 gives the following result, that is studied in more detail in Section 4.

Corollary 1.4 *There is a functor $\mathbf{cMon}(\mathbf{BP}(\mathbf{Set-Cat})) \rightarrow \mathbf{SMod-Cat}$.*

Proof. If \mathbf{C} and \mathbf{D} are symmetric monoidal categories, then a symmetric monoidal functor $\mathbf{C} \rightarrow \mathbf{D}$ restricts to a functor $\mathbf{cMon}(\mathbf{C}) \rightarrow \mathbf{cMon}(\mathbf{D})$. So Theorem 1.1 provides a functor $\mathbf{cMon}(\mathbf{BP}(\mathbf{Set-Cat})) \rightarrow \mathbf{cMon}(\mathbf{cMon-Cat})$. The general version of Theorem 1.3 that we will prove later provides a functor $\mathbf{cMon}(\mathbf{cMon-Cat}) \rightarrow \mathbf{Act}(\mathbf{cMon})\text{-Cat}$. Composing both functors establishes the result. \square

Summarizing, we have the following diagram.

$$\begin{array}{ccccc}
 & & (\mathbf{SMod}(\mathbf{Set}))\text{-Cat} & \longrightarrow & (\mathbf{cMon}(\mathbf{Set}))\text{-Cat} \\
 & \nearrow & \downarrow & & \downarrow \\
 \mathbf{cMon}(\mathbf{BP}(\mathbf{Set-Cat})) & \longrightarrow & \mathbf{BP}(\mathbf{Set-Cat}) & \longrightarrow & \mathbf{Cat} \\
 \downarrow & & \downarrow & & \downarrow \\
 & \nearrow & (\mathbf{Act}(\mathbf{Set}))\text{-Cat} & \longrightarrow & \mathbf{Cat} \\
 \mathbf{cMon}(\mathbf{Set-Cat}) & \longrightarrow & \mathbf{Set-Cat} & \longrightarrow & \mathbf{Cat}
 \end{array} \tag{5}$$

In order for the combination of Theorems 1.1 and 1.3 to make sense, the whole story, or at least Theorem 1.3, has to be generalized to an enriched setting. Sections 2–4 explain the above diagram in more detail, and generalize it to enriching categories \mathbf{V} that not necessarily equal \mathbf{Set} . Section 5 applies this to make precise the intuition that compact objects are “finite-dimensional” in these cases.

2 Additive enrichment

This section proves that a \mathbf{V} -category with finite biproducts can be seen as a $\mathbf{cMon}(\mathbf{V})$ -category, and that this change of perspective is functorial. Before this general version of Theorem 1.1 can even be stated, we need to consider an enriched version of the notion of biproduct. The following theorem characterizes algebraically how to recognize when a tensor product is in fact a coproduct, in a way reminiscent of [6] (see also [5]).

Theorem 2.1 *Let $(\mathbf{C}, \oplus, 0)$ be a symmetric monoidal category. Then $(\oplus, 0)$ provides finite coproducts if and only if the forgetful functor $\mathbf{cMon}(\mathbf{C}) \rightarrow \mathbf{C}$ is an isomorphism of categories.*

Proof. Suppose that $(\oplus, 0)$ provides finite coproducts, with the coherence maps α, λ and ρ induced by the coproducts. Let $U : \mathbf{cMon}(\mathbf{C}) \rightarrow \mathbf{C}$ be the forgetful

³ Objects of $\mathbf{cMon}(\mathbf{Cat})$ are *strict* symmetric monoidal categories, *i.e.* categories in which there are not just isomorphisms between $X \otimes (Y \otimes Z)$ and $(X \otimes Y) \otimes Z$ and so on, but equalities. To ease the presentation, we therefore mean *strict* (symmetric) monoidal categories whenever we talk about categories with tensor products, especially in the general \mathbf{V} -enriched setting, although some results are more general.

functor acting on objects as $U(X, \mu, \eta) = X$ and on morphisms as $U(f) = f$. Define $C : \mathbf{C} \rightarrow \mathbf{cMon}(\mathbf{C})$ on objects as $C(X) = (X, \nabla, u)$, where $\nabla = [\text{id}_X, \text{id}_X] : X \oplus X \rightarrow X$, and u is the unique morphism $0 \rightarrow X$. On a morphism f , it acts as $C(f) = f$. Then trivially $U \circ C = \text{Id}$. To prove that also $C \circ U = \text{Id}$, we show that there can be only one (commutative) monoid structure on $X \in \mathbf{C}$ with respect to $(\oplus, 0)$, *i.e.* for any $(X, \mu, \eta) \in \mathbf{cMon}(\mathbf{C})$ one has that $\mu = [\text{id}, \text{id}]$. This suffices because η is necessarily the unique morphism $0 \rightarrow X$. We have

$$\mu \circ \kappa_1 = \mu \circ [\kappa_1, \kappa_2 \circ u] \circ \kappa_1 = \mu \circ (\text{id} \oplus u) \circ \kappa_1 = \rho^{-1} \circ \kappa_1 = \text{id},$$

since $\kappa_1 : X \rightarrow X \oplus 0$ equals the coherence isomorphism ρ . Likewise $\mu \circ \kappa_2 = \text{id}$, so $\mu = [\text{id}, \text{id}]$, as needed.

Conversely, suppose that $\mathbf{cMon}(\mathbf{C}) \xrightleftharpoons[C]{U} \mathbf{C}$ is an isomorphism. By definition $U(X, \mu, \eta) = X$, so the monoid $C(X)$ is carried by X . Since C is a functor, the monoid structure maps, say $\nabla_X : X \oplus X \rightarrow X$ and $u_X : 0 \rightarrow X$, are natural in X . We first prove that 0 is an initial object. We have that $(0, \nabla_0, u_0)$ and $(0, \lambda_0, \text{id}_0)$ are both monoids (in \mathbf{C}). Moreover they satisfy the Hilton-Eckmann condition (A.1), so by Theorem A.1 in Appendix A we have $u_0 = \text{id}_0$. Naturality of u yields

$$f = f \circ \text{id}_0 = f \circ u_0 = u_X$$

for any $f : 0 \rightarrow X$. Hence u_X is the unique morphism $0 \rightarrow X$, and 0 is indeed an initial object. Finally, we show that $X \oplus Y$ is a coproduct of X and Y . Define as coproduct injections $\kappa : X \xrightarrow{\rho} X \oplus 0 \xrightarrow{\text{id} \oplus u_Y} X \oplus Y$ and $\kappa' : Y \xrightarrow{\lambda} 0 \oplus Y \xrightarrow{u_X \oplus \text{id}} X \oplus Y$. For given $f : X \rightarrow Z$ and $g : Y \rightarrow Z$, put $[f, g] = \nabla_Z \circ (f \oplus g) : X \oplus Y \rightarrow Z$. Then

$$\begin{aligned} [f, g] \circ \kappa_X &= \nabla_Z \circ (f \oplus g) \circ (\text{id} \oplus u_Y) \circ \rho \\ &= \nabla_Z \circ (\text{id} \oplus (g \circ u_Y)) \circ (f \oplus \text{id}) \circ \rho && (u \text{ natural}) \\ &= \nabla_Z \circ (\text{id} \oplus u_Y) \circ (f \oplus \text{id}) \circ \rho && (\nabla, u \text{ monoid}) \\ &= \rho^{-1} \circ (f \oplus \text{id}) \circ \rho = f. && (\rho \text{ natural}) \end{aligned}$$

Analogously, $[f, g] \circ \kappa_Y = g$. Moreover, $[f, g]$ is the unique such map since

$$\begin{aligned} [\kappa_X, \kappa_Y] &= (\nabla_X \oplus \nabla_Y) \circ (\text{id} \oplus \gamma \oplus \text{id}) \circ (\text{id} \oplus u_Y \oplus \text{id}) \\ &\quad \circ (\rho \oplus \text{id}) \circ (\text{id} \oplus u_X \oplus \text{id}) \circ (\text{id} \oplus \lambda) \\ &= (\nabla_X \oplus \text{id}) \circ (\text{id} \oplus \nabla_Y) \circ (\text{id} \oplus u_Y \oplus \text{id}) \circ (\text{id} \oplus \gamma \oplus \text{id}) \\ &\quad \circ (\rho \oplus \text{id}) \circ (\text{id} \oplus u_X \oplus \text{id}) \circ (\text{id} \oplus \lambda) \\ &= (\nabla_X \oplus \text{id}) \circ (\text{id} \oplus \lambda^{-1}) \circ (\text{id} \oplus \gamma \oplus \text{id}) \\ &\quad \circ (\rho \oplus \text{id}) \circ (\text{id} \oplus u_X \oplus \text{id}) \circ (\text{id} \oplus \lambda) \\ &= (\nabla_X \oplus \text{id}) \circ (\text{id} \oplus u_X \oplus \text{id}) \circ (\text{id} \oplus \lambda) \\ &= (\text{id} \oplus \lambda^{-1}) \circ (\text{id} \oplus \lambda) = \text{id}. \end{aligned}$$

□

Dually, a symmetric monoidal structure $(\oplus, 0)$ on a category \mathbf{C} provides finite products if and only if \mathbf{C}^{op} is isomorphic to the category of commutative comonoids

in \mathbf{C} . Let us write $\mathbf{C}^{\rightleftarrows}$ for the full subcategory of $\mathbf{C}^{\text{op}} \times \mathbf{C}$ consisting of objects (X, X) . Its objects can be identified with those of \mathbf{C} , and a morphism $f : X \rightarrow Y$ in $\mathbf{C}^{\rightleftarrows}$ is a pair of morphisms $f^{\leftarrow} : Y \rightarrow X$, $f^{\rightarrow} : X \rightarrow Y$ of \mathbf{C} . If \mathbf{C} is (symmetric) monoidal (closed), so is $\mathbf{C}^{\rightleftarrows}$. Hence an object in $\mathbf{cMon}(\mathbf{C}^{\rightleftarrows})$ consists of a monoid and a comonoid in \mathbf{C} with the same carrier object. Thus we arrive at the following algebraic characterization of how to recognize when a tensor product in fact provides finite biproducts.

Corollary 2.2 *Let $(\mathbf{C}, \oplus, 0)$ be a symmetric monoidal category. Then $(\oplus, 0)$ provides finite biproducts if and only if the underlying functor $\mathbf{cMon}(\mathbf{C}^{\rightleftarrows}) \rightarrow \mathbf{C}^{\rightleftarrows}$ is an isomorphism.*

Proof. Suppose the underlying functor is an isomorphism. Then 0 is an initial object by Theorem 2.1, and a terminal object by the dual of that Theorem. Hence 0 is a zero object. Condition (1)–(3) are satisfied; let us show *e.g.* condition (2):

$$\pi_X \circ \kappa_X = \rho^{-1} \circ (\text{id} \oplus n_Y) \circ (\text{id} \oplus u_Y) \circ \rho = \rho^{-1} \circ (\text{id} \oplus u_Y) \circ \rho = \text{id}_X.$$

The converse is trivial: if \mathbf{C} has coinciding finite products and coproducts, then Theorem 2.1 and its dual show that the underlying functor is an isomorphism. (See also footnote 3.) \square

Compare this to the following: a category \mathbf{C} is self-dual, *i.e.* $\mathbf{C} \cong \mathbf{C}^{\text{op}}$, if and only if the forgetful functor $\mathbf{C}^{\rightleftarrows} \rightarrow \mathbf{C}$ is an isomorphism.

We now turn this algebraic characterization of biproducts in locally small categories into a definition of biproducts in \mathbf{V} -categories. The notion of monoid in a monoidal category duly enriches. For an enriching (symmetric) monoidal category \mathbf{V} , we speak of $\mathbf{C} \in \mathbf{cMon}(\mathbf{V}\text{-Cat})$ as a ‘(symmetric) monoidal \mathbf{V} -category’. This means that \mathbf{C} comes equipped with morphisms

$$\otimes_{\mathbf{C}} : \mathbf{C}(X, Y) \otimes_{\mathbf{V}} \mathbf{C}(X', Y') \rightarrow \mathbf{C}(X \otimes_{\mathbf{C}} X', Y \otimes_{\mathbf{C}} Y')$$

in \mathbf{V} . A (commutative) monoid object in \mathbf{C} then consists of an object $X \in \mathbf{C}$ together with morphisms $\mu : I_{\mathbf{V}} \rightarrow \mathbf{C}(X \otimes_{\mathbf{C}} X, X)$ and $\eta : I_{\mathbf{V}} \rightarrow \mathbf{C}(I_{\mathbf{C}}, X)$, making the appropriate diagrams (in \mathbf{V}) commute. The definition of $\mathbf{C}^{\rightleftarrows}$ above also works for \mathbf{V} -enriched categories \mathbf{C} . Hence Corollary 2.2 enables us to talk about finite biproducts in \mathbf{V} -enriched categories without having to resort to product structure on \mathbf{V} or weighted (co)limits. In particular, it is more general than the usual notion of a \mathbf{V} -coproduct, which needs \mathbf{V} to have finite products — if \mathbf{V} is a category with finite products, then finite coproduct structure on a \mathbf{V} -category \mathbf{C} is traditionally regarded as \mathbf{V} -natural isomorphisms $\mathbf{C}(X \oplus Y, -) \cong \mathbf{C}(X, -) \times \mathbf{C}(Y, -)$ and $\mathbf{C}(0, -) \cong 1$ [11]. Let us collect all enriched categories with finite biproducts in a category. For a symmetric monoidal category \mathbf{V} , define $\mathbf{BP}(\mathbf{V}\text{-Cat})$ as the full subcategory of $\mathbf{cMon}(\mathbf{V}\text{-Cat})$ consisting of \mathbf{V} -categories \mathbf{C} such that the underlying functor $\mathbf{cMon}(\mathbf{C}^{\rightleftarrows}) \rightarrow \mathbf{C}^{\rightleftarrows}$ is an isomorphism.

We are now in a position to tackle the general version of Theorem 1.1. The next theorem proves that \mathbf{V} -enrichment of a category that has biproducts in the above sense can be lifted to an enrichment over $\mathbf{cMon}(\mathbf{V})$.

Theorem 2.3 *Let \mathbf{V} be a symmetric monoidal category. There is a functor*

$$(-)^\oplus : \mathbf{BP}(\mathbf{V}\text{-Cat}) \longrightarrow (\mathbf{cMon}(\mathbf{V}))\text{-Cat}.$$

Proof. Let $(\mathbf{C}, \oplus, 0) \in \mathbf{BP}(\mathbf{V}\text{-Cat})$. It comes equipped with $\Delta_X : I_{\mathbf{V}} \rightarrow \mathbf{C}(X, X \oplus X)$, $\nabla_X : I_{\mathbf{V}} \rightarrow \mathbf{C}(X \oplus X, X)$, $u_X : I_{\mathbf{V}} \rightarrow \mathbf{C}(0, X)$ and $n_X : I_{\mathbf{V}} \rightarrow \mathbf{C}(X, 0)$. The objects of \mathbf{C}^\oplus are those of \mathbf{C} . The carrier of the homobject $\mathbf{C}^\oplus(X, Y)$ is $\mathbf{C}(X, Y) \in \mathbf{V}$. Its monoid unit $0_{XY} : I_{\mathbf{V}} \rightarrow \mathbf{C}(X, Y)$ is given by

$$0_{XY} : I_{\mathbf{V}} \xrightarrow{\cong} I_{\mathbf{V}} \otimes I_{\mathbf{V}} \xrightarrow{n_X \otimes u_Y} \mathbf{C}(X, 0) \otimes \mathbf{C}(0, Y) \xrightarrow{\circ_{\mathbf{C}}} \mathbf{C}(X, Y)$$

The monoid multiplication $+$: $\mathbf{C}(X, Y) \otimes \mathbf{C}(X, Y) \rightarrow \mathbf{C}(X, Y)$ of the homobject $\underline{\mathbf{C}}(X, Y)$ is given as follows.

$$\begin{array}{ccc} \mathbf{C}(X, Y) \otimes \mathbf{C}(X, Y) & \xrightarrow{\quad + \quad} & \mathbf{C}(X, Y) \\ \oplus \downarrow & & \uparrow (\circ_{\mathbf{C}}) \circ (\text{id} \otimes \circ_{\mathbf{C}}) \\ \mathbf{C}(X \oplus X, Y \oplus Y) & \xrightarrow{\quad \Delta_X \otimes \text{id} \otimes \nabla_Y \quad} & \mathbf{C}(X \oplus X, Y \oplus Y) \otimes \mathbf{C}(Y \oplus Y, Y) \\ \cong \downarrow & & \uparrow \\ I_{\mathbf{V}} \otimes \mathbf{C}(X \oplus X, Y \oplus Y) & \otimes I_{\mathbf{V}} & \end{array}$$

Composition $\circ_{\mathbf{C}}$ is a monoid morphism for this structure because \oplus is a \mathbf{V} -functor. We leave it to the reader to show that this data indeed defines a commutative monoid; essentially it is an enriched version of the argument for **Set**-categories.

Since a morphism in $\mathbf{BP}(\mathbf{V}\text{-Cat})$ is a \mathbf{V} -functor that (strictly) preserves the biproduct structure, and because Δ, ∇, n, u are natural, this assignment $\mathbf{C} \mapsto \mathbf{C}^\oplus$ is functorial. \square

3 Scalar multiplication enrichment

The introduction set the first step towards proving that every monoidal category \mathbf{C} is enriched over $\mathbf{Act}_{\mathbf{C}(I, I)}(\mathbf{Set})$. Before we consider the general enriched situation, let us complete the reasoning in the **Set**-enriched setting, by ensuring that composition is a morphism in $\mathbf{Act}_{\mathbf{C}(I, I)}(\mathbf{Set})$. Because $s \bullet (g \circ f) = g \circ (s \bullet f) = (s \bullet g) \circ f$, this requires that $\mathbf{Act}_{\mathbf{C}(I, I)}(\mathbf{Set})$ is equipped with a tensor structure that is universal for bimorphisms. This will occupy us for the next few lemmas, that apply techniques developed by Kock and Day ([13], but see also [10]).

Lemma 3.1 *If \mathbf{C} is a locally small (symmetric) monoidal category, then $\mathbf{C}(I, I)$ is a (commutative) monoid.*

Proof. [12,1] By the Hilton-Eckmann argument (see Appendix A) the monoid $(\mathbf{C}(I, I), \bullet, \text{id}_I)$ coincides with $(\mathbf{C}(I, I), \circ, \text{id}_I)$, and is in fact commutative. The

following diagram establishes commutativity directly.

$$\begin{array}{ccccc}
 I & \xrightarrow[\lambda]{\cong} & I \otimes I & \xlongequal{\quad} & I \otimes I & \xrightarrow[\rho^{-1}]{\cong} & I \\
 \uparrow s & & \uparrow s \otimes \text{id} & & \downarrow \text{id} \otimes t & & \downarrow t \\
 I & \xrightarrow[\lambda=\rho]{\cong} & I \otimes I & \xrightarrow{s \otimes t} & I \otimes I & \xrightarrow[\lambda^{-1}=\rho^{-1}]{\cong} & I \\
 \downarrow t & & \downarrow \text{id} \otimes t & & \uparrow s \otimes \text{id} & & \uparrow s \\
 I & \xrightarrow[\rho]{\cong} & I \otimes I & \xlongequal{\quad} & I \otimes I & \xrightarrow[\lambda^{-1}]{\cong} & I
 \end{array}$$

Notice how this essentially uses the coherence axiom $\lambda_I = \rho_I$. \square

If \mathbf{V} is a symmetric monoidal category, then so is $\mathbf{V}\text{-Cat}$ [3]. Hence it makes sense to talk of (strict) monoidal \mathbf{V} -categories as objects of $\mathbf{Mon}(\mathbf{V}\text{-Cat})$. In fact, the previous lemma extends functorially and enriches: if \mathbf{V} is a symmetric monoidal category, then there is a functor $\mathbf{cMon}(\mathbf{V}\text{-Cat}) \rightarrow \mathbf{cMon}(\mathbf{V})$.

Recall that a monad T on a symmetric monoidal category is called *strong* if there is a “strength” natural transformation $st : X \otimes TY \rightarrow T(X \otimes Y)$ satisfying certain conditions. The monad is called *commutative* if both iterated “double strength” maps $TX \otimes TY \rightarrow T(X \otimes Y)$ coincide [10, Definition 3.4].

Lemma 3.2 *If \mathbf{V} is a monoidal category and $M \in \mathbf{Mon}(\mathbf{V})$, then $M \otimes (-) : \mathbf{V} \rightarrow \mathbf{V}$ is a monad, whose category of algebras is $\mathbf{Act}_M(\mathbf{V})$. If \mathbf{V} is symmetric monoidal, then the monad $M \otimes (-)$ is strong. The monad $M \otimes (-)$ is commutative if and only if the monoid M is.*

Proof. The unit and multiplication of the monad are given by

$$\begin{aligned}
 \eta : X &\xrightarrow{\cong} I \otimes X \xrightarrow{e \otimes \text{id}} M \otimes X, \\
 \mu : M \otimes (M \otimes X) &\xrightarrow{\cong} (M \otimes M) \otimes X \xrightarrow{m \otimes \text{id}} M \otimes X.
 \end{aligned}$$

If \mathbf{C} is symmetric monoidal, then there is a strength map

$$st : X \otimes (M \otimes Y) \cong (X \otimes M) \otimes Y \xrightarrow{\gamma \otimes \text{id}} (M \otimes X) \otimes Y \cong M \otimes (X \otimes Y).$$

The double strength maps boil down to

$$\begin{array}{ccc}
 (M \otimes X) \otimes (M \otimes Y) & \xrightarrow{\text{dst}} & M \otimes (X \otimes Y) \\
 \parallel & & \parallel \\
 (M \otimes X) \otimes (M \otimes Y) & \xrightarrow{\cong} (M \otimes M) \otimes (X \otimes Y) \xrightarrow{m \otimes \text{id}} & M \otimes (X \otimes Y) \\
 \parallel & \downarrow \gamma \otimes \text{id} & \\
 & (M \otimes M) \otimes (X \otimes Y) \xrightarrow{m \otimes \text{id}} & M \otimes (X \otimes Y) \\
 \parallel & & \parallel \\
 (M \otimes X) \otimes (M \otimes Y) & \xrightarrow{\text{dst}'} & M \otimes (X \otimes Y)
 \end{array}$$

Hence they coincide if and only if the monoid M is commutative. \square

Lemma 3.3 *Let \mathbf{V} be a symmetric monoidal closed category, $M \in \mathbf{cMon}(\mathbf{V})$, and suppose $\mathbf{Act}_M(\mathbf{V})$ has coequalizers of reflexive pairs. Then $\mathbf{Act}_M(\mathbf{V})$ is a symmetric monoidal closed category, in which $X \otimes Y$ is universal such that every bimorphism $X \times Y \rightarrow Z$ factors through it.*

Proof. Apply [10, Lemma 5.3]. The resulting tensor product structure is such that there are universal bimorphisms [10, Lemma 5.1]. \square

A special case of the previous lemma is $\mathbf{V} = \mathbf{Set}$. In this case $\mathbf{Act}_M(\mathbf{V})$ has coequalizers of reflexive pairs since it is in fact a topos. For $X, Y \in \mathbf{Act}_M(\mathbf{Set})$, the tensor product $X \otimes Y$ is given explicitly by $X \times Y / \sim$, where \sim is the least equivalence relation determined by $(m \cdot x, y) \sim (x, m \cdot y)$, with action given by $m \cdot [x, y] = [m \cdot x, y] = [x, m \cdot y]$.

There is more structure behind the previous lemma than stated there. In fact, since $V \mapsto M \otimes V$ is the free functor $\mathbf{V} \rightarrow \mathbf{Act}_M(\mathbf{V})$, it is a morphism of symmetric monoidal categories, *i.e.* it preserves the symmetric monoidal structure⁴. In particular, there is an isomorphism $I_{\mathbf{Act}_M(\mathbf{V})} \cong M \otimes I_{\mathbf{V}}$ of (commutative) monoids. In case $\mathbf{V} = \mathbf{Set}$ and $M = \mathbf{C}(I, I)$ this shows that the monoidal structure on \mathbf{V} provided by Lemma 3.3 and that on \mathbf{C} have ‘the same scalars’.

Note 1 *From now on, we will assume that the symmetric monoidal category \mathbf{V} is such that $\mathbf{Act}_{\mathbf{C}(I, I)}(\mathbf{V})$ has coequalizers of reflexive pairs. This ensures that $\mathbf{Act}_{\mathbf{C}(I, I)}(\mathbf{V})$ is again a symmetric monoidal category, and hence it makes sense to talk about enrichment over it. We will only use this technical assumption for this reason.*

A common scenario in which the above assumption is fulfilled is when \mathbf{V} is a symmetric monoidal closed category that has coequalizers of reflexive pairs, like in the situation $\mathbf{V} = \mathbf{Set}$ above. (In fact, it suffices if $\mathbf{C}(I, I) \otimes (-)$ has a right adjoint.)

Now that we have developed a monoidal structure on $\mathbf{Act}_{\mathbf{C}(I, I)}(\mathbf{Set})$ that is universal for bimorphisms, we can use the scalar multiplication action as an enrichment. The following theorem summarizes this main insight.

Theorem 3.4 *Every locally small symmetric monoidal category \mathbf{C} is enriched over the symmetric monoidal closed category $\mathbf{V} = \mathbf{Act}_{\mathbf{C}(I, I)}(\mathbf{Set})$, in such a way that $I_{\mathbf{V}} \cong \mathbf{C}(I, I)$.*

Proof. Put the homobject $\mathbf{C}(X, Y) \in \mathbf{Act}_{\mathbf{C}(I, I)}(\mathbf{Set})$ to be the set $\mathbf{C}(X, Y)$ with the action $\alpha : \mathbf{C}(X, Y) \times \mathbf{C}(I, I) \rightarrow \mathbf{C}(X, Y)$ on it given by scalar multiplication as $\alpha(f, s) = s \bullet f$. The composition morphism $\mathbf{C}(X, Y) \otimes \mathbf{C}(Y, Z) \rightarrow \mathbf{C}(X, Z)$ is now the unique one through which the bimorphism $\mathbf{C}(X, Y) \times \mathbf{C}(Y, Z) \rightarrow \mathbf{C}(X, Z)$, determined by $(f, g) \mapsto g \circ f$, factors; this is a morphism in $\mathbf{Act}_{\mathbf{C}(I, I)}(\mathbf{Set})$. The identity morphism $1 \rightarrow \mathbf{C}(X, X)$ is given by $*$ $\mapsto \text{id}_X$; this is also a morphism in $\mathbf{Act}_{\mathbf{C}(I, I)}(\mathbf{Set})$ since 1 carries the trivial action. One easily verifies that these satisfy the requirements of an enriched category. \square

Whereas the previous theorem covers the case of \mathbf{Set} -enriched categories, the following theorem gives the general construction. It incorporates functoriality; but for that we first need to get rid of the indexing monoid M in $\mathbf{Act}_M(\mathbf{V})$. We denote by $\mathbf{Act}(\mathbf{V})$ the Grothendieck completion $\int_{M \in \mathbf{cMon}(\mathbf{V})} \mathbf{Act}_M(\mathbf{V})$ of the indexed category $\mathbf{cMon}(\mathbf{V})^{\text{op}} \rightarrow \mathbf{Cat}$. Explicitly, $\mathbf{Act}(\mathbf{V})$ has as objects pairs (M, α) with $M \in \mathbf{cMon}(\mathbf{V})$ and $\alpha \in \mathbf{Act}_M(\mathbf{V})$. Morphisms from $(M, \alpha : M \otimes X \rightarrow X)$ to $(N, \beta : N \otimes Y \rightarrow Y)$ in $\mathbf{Act}(\mathbf{V})$ are pairs of morphisms $f : M \rightarrow N$ and $g : X \rightarrow Y$

⁴ If moreover \mathbf{V} is a topos, it is also a geometric morphism of toposes [16].

satisfying $\beta \circ (f \otimes g) = g \circ \alpha$ [8]. If \mathbf{V} is symmetric monoidal, then so is $\mathbf{Act}(\mathbf{V})$, whence it makes sense to talk of $(\mathbf{Act}(\mathbf{V}))$ -enriched categories.

Theorem 3.5 *Let \mathbf{V} be a symmetric monoidal category. There is a functor*

$$(-)^\otimes : \mathbf{cMon}(\mathbf{V}\text{-Cat}) \longrightarrow (\mathbf{Act}(\mathbf{V}))\text{-Cat}.$$

Proof. We first describe how $(-)^{\otimes}$ works on objects. Let $\mathbf{C} \in \mathbf{cMon}(\mathbf{V}\text{-Cat})$. This means that \mathbf{C} is a \mathbf{V} -enriched category, and hence comes equipped with *e.g.* morphisms $i_X : I_{\mathbf{V}} \rightarrow \mathbf{C}(X, X)$. Furthermore, it means that there is a \mathbf{V} -functor $\otimes_{\mathbf{C}}$; explicitly, we are given a morphism $\otimes_{\mathbf{C}} : |\mathbf{C}| \times |\mathbf{C}| \rightarrow |\mathbf{C}|$ in \mathbf{Set} , and a morphism $\otimes_{\mathbf{C}} : \mathbf{C}(X, X') \otimes_{\mathbf{V}} \mathbf{C}(Y, Y') \rightarrow \mathbf{C}(X \otimes_{\mathbf{C}} X', Y \otimes_{\mathbf{C}} Y')$ in \mathbf{V} . Finally, it means we are given an object $I_{\mathbf{C}} \in |\mathbf{C}|$. These data satisfy the (strict) monoid requirements, like $I_{\mathbf{C}} \otimes_{\mathbf{C}} X = X$.

The objects of the \mathbf{V} -enriched category \mathbf{C} and the $\mathbf{Act}(\mathbf{V})$ -enriched category \mathbf{C}^\otimes are the same: $|\mathbf{C}^\otimes| = |\mathbf{C}|$. The homobjects are determined by the action of scalar multiplication, *i.e.* $\mathbf{C}^\otimes(X, Y)$ is the action

$$\mathbf{C}(I, I) \otimes_{\mathbf{V}} \mathbf{C}(X, Y) \xrightarrow{\otimes_{\mathbf{C}}} \mathbf{C}(I \otimes_{\mathbf{C}} X, I \otimes_{\mathbf{C}} Y) = \mathbf{C}(X, Y).$$

The identity on X is the morphism

$$\mathbf{C}(I, I) \xrightarrow{\cong} \mathbf{C}(I, I) \otimes_{\mathbf{V}} I_{\mathbf{V}} \xrightarrow{\text{id} \otimes i_X} \mathbf{C}(I, I) \otimes \mathbf{C}(X, X) \xrightarrow{\otimes_{\mathbf{C}}} \mathbf{C}(X, X)$$

in \mathbf{V} . It is a morphism of actions $I_{\mathbf{Act}_{\mathbf{C}(I, I)}(\mathbf{V})} \rightarrow \mathbf{C}^\otimes(X, X)$ since $I_{\mathbf{Act}_M(\mathbf{V})}$ is the action $M \otimes M \xrightarrow{m} M$. An involved but straightforward calculation, that uses the structure of the tensor product in $\mathbf{Act}_M(\mathbf{V})$, now shows that these data in fact provide an enrichment.

We now turn to the action of $(-)^{\otimes}$ on morphisms. Let F be a morphism $\mathbf{C} \rightarrow \mathbf{D}$ in $\mathbf{cMon}(\mathbf{V}\text{-Cat})$. Define its image F^\otimes to work on objects $X \in \mathbf{C}^\otimes$ as $F^\otimes(X) = F(X)$. It also works on morphisms as F — since F is a ‘(strict) monoidal \mathbf{V} -functor’, it is automatically a (scalar multiplication) action morphism.

$$\begin{array}{ccc} \mathbf{C}(I, I) \otimes_{\mathbf{V}} \mathbf{C}(X, Y) & \xrightarrow{\otimes_{\mathbf{C}}} & \mathbf{C}(I \otimes_{\mathbf{C}} X, I \otimes_{\mathbf{C}} Y) = \mathbf{C}(X, Y) \\ F_{II} \otimes_{\mathbf{V}} F_{XY} \downarrow & & \downarrow F_{XY} \quad \downarrow F_{XY} \\ \mathbf{D}(I, I) \otimes_{\mathbf{V}} \mathbf{D}(FX, FY) & \xrightarrow{\otimes_{\mathbf{D}}} & \mathbf{D}(I \otimes_{\mathbf{D}} FX, I \otimes_{\mathbf{D}} FY) = \mathbf{D}(FX, FY) \end{array}$$

That is, F_{XY}^\otimes is indeed a morphism in $\mathbf{Act}(\mathbf{V})$. □

The extension of enrichment in the previous theorem is initial, in the sense that the ‘forgetful’ functor $\mathbf{C}(I, -) : \mathbf{C} \rightarrow \mathbf{V}$ of any symmetric monoidal \mathbf{V} -enriched category \mathbf{C} factors through it, as in the following commutative diagram of monoidal functors.

$$\begin{array}{ccc} \mathbf{C} & \xrightarrow{\mathbf{C}(I, -)} & \mathbf{V} \\ & \searrow \mathbf{C}(I, -) & \nearrow (\alpha : \mathbf{C}(I, I) \otimes_{\mathbf{V}} \mathbf{V} \rightarrow \mathbf{V}) \\ & \mathbf{Act}_{\mathbf{C}(I, I)}(\mathbf{V}) & \end{array}$$

4 Semimodule enrichment

This section considers the situation arising from a monoidal \mathbf{V} -category with finite biproducts. By the previous two sections, we know that such a category can be seen both as a $\mathbf{cMon}(\mathbf{V})$ -category and as an $\mathbf{Act}(\mathbf{V})$ -category. Using the correct distributivity we will show that it can in fact be seen as a category enriched over semimodules.

When we say ‘a monoidal \mathbf{V} -category \mathbf{C} with finite biproducts’, we mean $\mathbf{C} \in \mathbf{cMon}(\mathbf{BP}(\mathbf{V}\text{-Cat}))$.⁵ After all, even if $\mathbf{V} = \mathbf{Set}$ the tensor product distributes over the biproduct by

$$\begin{aligned} \tau &= \langle \text{id} \otimes \pi_1, \text{id} \otimes \pi_2 \rangle : X \otimes (Y \oplus Z) \xrightarrow{\cong} (X \otimes Y) \oplus (X \otimes Z), \\ \tau^{-1} &= [\text{id} \otimes \kappa_1, \text{id} \otimes \kappa_2], \end{aligned}$$

but the biproduct rarely distributes over the tensor product, so that considering $\mathbf{BP}(\mathbf{cMon}(\mathbf{V}\text{-Cat}))$ is not so sensible. This resonates well with the observation in the introduction that $\mathbf{cSRng}(\mathbf{Set})$ equals $\mathbf{cMon}(\mathbf{cMon}(\mathbf{Set}))$, where the multiplicative structure distributes over the additive structure. The stock example is the category of complex vector spaces and linear functions, in which one can multiply (tensor) and add (direct sum) objects as well as morphisms.

A first sign of this distributivity is seen when reconsidering Lemma 3.1 in a monoidal \mathbf{V} -category that moreover has finite biproducts. Then $\mathbf{C}(I, I)$ is not just a monoid, but in fact a commutative semiring.

Corollary 4.1 *Let \mathbf{V} be a symmetric monoidal category. There is a functor*

$$\mathbf{cMon}(\mathbf{BP}(\mathbf{V}\text{-Cat})) \longrightarrow \mathbf{cSRng}(\mathbf{V}).$$

Proof. It acts on objects by sending $(\mathbf{C}, \otimes, I, \oplus, 0)$ to $\mathbf{C}(I, I)$. We already saw that $(\mathbf{C}(I, I), \bullet, \text{id}_I)$ is a monoid, and that $(\mathbf{C}(I, I), +, 0)$ is a commutative monoid. Hence it suffices to verify distributivity and annihilation. We only consider the case $\mathbf{V} = \mathbf{Set}$, which enriches easily, so that the diagrams are clearer. Distributivity is established by the following commutative diagram.

$$\begin{array}{ccccccc} I & \xrightarrow{\hspace{10em}} & I & & I & & I \\ \parallel & & \parallel & & \parallel & & \parallel \\ I & \xrightarrow{\lambda} I \otimes I \xrightarrow{\text{id} \otimes \Delta} I \otimes (I \oplus I) \xrightarrow{s \otimes (t \oplus r)} I \otimes (I \oplus I) \xrightarrow{\text{id} \otimes \nabla} I \otimes I \xrightarrow{\lambda^{-1}} I & & & & & I \\ \parallel & & \parallel & & \parallel & & \parallel \\ I & \xrightarrow{\Delta} I \oplus I \xrightarrow{\lambda \oplus \lambda} (I \otimes I) \oplus (I \otimes I) \xrightarrow{(s \otimes t) \oplus (s \otimes r)} (I \otimes I) \oplus (I \otimes I) \xrightarrow{\lambda^{-1} \oplus \lambda^{-1}} I \oplus I \xrightarrow{\nabla} I & & & & & I \\ \parallel & & \parallel & & \parallel & & \parallel \\ I & \xrightarrow{\hspace{10em}} & I & & I & & I \\ & & & & & & (s \bullet t) + (s \bullet r) \end{array}$$

⁵ This automatically takes care of coherence in the not-necessarily strict case [14]. See also footnote 3.

Annihilation is shown as follows.

$$\begin{array}{ccccc}
 I & \xrightarrow{\quad \quad \quad \dagger \quad \quad \quad} & & & 0 \\
 \parallel & & & & \parallel \\
 I & \xrightarrow{\lambda} I \otimes I \xrightarrow{\dagger \otimes \text{id}} 0 \otimes I \xrightarrow{\rho^{-1}} & & & 0 \\
 \bullet \circ s \downarrow & & 0 \otimes s \downarrow & & \downarrow \text{id} \otimes s \\
 I & \xleftarrow{\lambda^{-1}} I \otimes I \xleftarrow{\dagger \otimes \text{id}} 0 \otimes I \xleftarrow{\rho^{-1}} & & & 0 \\
 \parallel & & & & \parallel \\
 I & \xleftarrow{\quad \quad \quad \dagger \quad \quad \quad} & & & 0
 \end{array}$$

So indeed $\mathbf{C}(I, I) \in \mathbf{cSRng}(\mathbf{V})$. Because a morphism in $\mathbf{cMon}(\mathbf{BP}(\mathbf{V}\text{-Cat}))$ is a \mathbf{V} -functor that preserves both the monoidal structure and the biproduct, and a morphism in $\mathbf{cSRng}(\mathbf{V}) = \mathbf{cMon}(\mathbf{cMon}(\mathbf{V}))$ is a morphism that respects $\bullet, 1, +$ and 0 , the above assignment is functorial. \square

Let us denote by $\mathbf{SMod}(\mathbf{V}) = \mathbf{Act}(\mathbf{cMon}(\mathbf{V}))$ the category of semimodules over an arbitrary semiring in \mathbf{V} . (In particular, $\mathbf{SMod}(\mathbf{V})$ is symmetric monoidal; in case $\mathbf{V} = \mathbf{Set}$, the induced monoidal structure coincides with that given by extension-of-scalars [17].)

Corollary 4.2 *Let \mathbf{V} be a symmetric monoidal category. There is a functor*

$$(-)^{\oplus \otimes} : \mathbf{cMon}(\mathbf{BP}(\mathbf{V}\text{-Cat})) \longrightarrow (\mathbf{SMod}(\mathbf{V}))\text{-Cat}.$$

Proof. Notice that $(-)^{\oplus}$ is a symmetric monoidal functor. That is, if $\mathbf{C}, \mathbf{D} \in \mathbf{cMon}(\mathbf{BP}(\mathbf{V}\text{-Cat}))$, then $(\mathbf{C} \otimes \mathbf{D})^{\oplus} \cong \mathbf{C}^{\oplus} \otimes \mathbf{D}^{\oplus}$, where the tensor product is that of \mathbf{V} -categories. Moreover, $\mathbf{1}^{\oplus} = \mathbf{1}$ [11,3]. Hence $(-)^{\oplus}$ restricts to $\mathbf{cMon}((-)^{\oplus}) : \mathbf{cMon}(\mathbf{BP}(\mathbf{V}\text{-Cat})) \rightarrow \mathbf{cMon}(\mathbf{cMon}(\mathbf{V}))\text{-Cat}$. The desired functor is then the composition $(-)^{\oplus \otimes} = (-)^{\otimes} \circ \mathbf{cMon}((-)^{\oplus})$. \square

If we combine all the functors so far with forgetful ones, we can summarize the entire article (so far) in the following diagram, generalizing diagram (5) from \mathbf{Set} to \mathbf{V} .

$$\begin{array}{ccccc}
 & & (\mathbf{SMod}(\mathbf{V}))\text{-Cat} & \longrightarrow & (\mathbf{cMon}(\mathbf{V}))\text{-Cat} \\
 & \nearrow & \downarrow & & \downarrow \\
 \mathbf{cMon}(\mathbf{BP}(\mathbf{V}\text{-Cat})) & \longrightarrow & \mathbf{BP}(\mathbf{V}\text{-Cat}) & \longrightarrow & \mathbf{BP}(\mathbf{V}\text{-Cat}) \\
 \downarrow & & \downarrow & & \downarrow \\
 & \nearrow & (\mathbf{Act}(\mathbf{V}))\text{-Cat} & \longrightarrow & \mathbf{Cat} \\
 \mathbf{cMon}(\mathbf{V}\text{-Cat}) & \longrightarrow & \mathbf{V}\text{-Cat} & \longrightarrow & \mathbf{V}\text{-Cat}
 \end{array}$$

5 Compactness and dimension

Categories with tensor products and biproducts have lately been studied as semantical models of quantum computation [2]. Especially compact objects in such categories play an important role. An object X of a symmetric monoidal category \mathbf{C} is called *compact* when there are an object $X^* \in \mathbf{C}$ and morphisms $\eta : I \rightarrow X^* \otimes X$

and $\varepsilon : X \otimes X^* \rightarrow I$ that make the following diagrams commute.

$$\begin{array}{ccc}
 X \xrightarrow{\cong} X \otimes I \xrightarrow{\text{id} \otimes \eta} X \otimes (X^* \otimes X) & & X^* \xrightarrow{\cong} I \otimes X^* \xrightarrow{\eta \otimes \text{id}} (X^* \otimes X) \otimes X^* \\
 \parallel & & \parallel \\
 X \xleftarrow{\cong} I \otimes X \xleftarrow{\varepsilon \otimes \text{id}} (X \otimes X^*) \otimes X & & X^* \xleftarrow{\cong} X^* \otimes I \xleftarrow{\text{id} \otimes \eta} X^* \otimes (X \otimes X^*)
 \end{array} \quad (6)$$

Compact objects ‘behave finite-dimensionally’. Indeed, in the prime example categories of complex vector spaces or Hilbert spaces, the compact objects are precisely the finite-dimensional ones. Our semimodule enrichment now puts us in a position to make this intuition precise.

Recall that an S -semimodule is called *free* if it is of the form S^X for some $X \in \mathbf{Set}$, with pointwise operations; the cardinality of X is its *dimension*. A semimodule is called *projective* if it is a retract of a free one [7]. If it is a retract of a finite-dimensional free semimodule, it is *projective of finite type*.

Lemma 5.1 *Let $S \in \mathbf{SRng}$. If M is a compact object in \mathbf{SMod}_S , then it is projective of finite type.*

Proof. [9, Lemma 1.3] \mathbf{SMod}_S is closed and hence enriched over itself. So if M is compact, then there is an isomorphism $a : \mathbf{SMod}_S(M, M) \rightarrow M \otimes \mathbf{SMod}_S(M, S)$ in \mathbf{SMod}_S . Say $a(\text{id}_M) = \sum_{i=1}^n m_i \otimes \varphi^i$ for $m_i \in M$ and $\varphi^i : M \rightarrow S$ and $i = 1, \dots, n$. Define $f : M \rightarrow S^n$ by $f(m) = (\varphi^1(m), \dots, \varphi^n(m))$, and $g : S^n \rightarrow M$ by $g(r_1, \dots, r_n) = \sum_{i=1}^n m_i r_i$. Diagram (6) then yields that $\sum_{i=1}^n m_i \varphi^i(m) = m$ for all $m \in M$, whence $g \circ f = \text{id}_M$. \square

Let \mathbf{C} be a locally small symmetric monoidal category with finite biproducts. One easily shows that the ‘forgetful’ functor $\mathbf{C}(I, -) : \mathbf{C} \rightarrow \mathbf{SMod}_{\mathbf{C}(I, I)}$ is monoidal, that is, there is a natural transformation with components $\mathbf{C}(I, X) \otimes \mathbf{C}(I, Y) \rightarrow \mathbf{C}(I, X \otimes Y)$. If it is moreover a natural isomorphism, it preserves compact objects. This is the case when the tensor product resembles the usual algebraic one used in standard quantum physics to model entanglement. Combining this with the previous lemma gives us the following precise version of the intuition that compact objects are ‘finite-dimensional’ in the standard setting.

Corollary 5.2 *If X is a compact object in a locally small symmetric monoidal category \mathbf{C} with finite biproducts, and $\mathbf{C}(I, -) : \mathbf{C} \rightarrow \mathbf{SMod}_{\mathbf{C}(I, I)}$ is strong monoidal, then $\mathbf{C}(I, X)$ is a projective semimodule of finite type.*

6 Future work

The fact that symmetric monoidal categories with biproducts are enriched over semimodules is a first step towards a representation theorem for categories that are Abelian except for an absence of subtraction. We intend to formulate the properties required of such categories to replace the well-behavedness of mono’s, kernels, epi’s and cokernels, and subsequently prove that every such category embeds into a category of semimodules. An interesting matter then is the relationship between Mitchell’s scalars and Abramsky’s scalars. Finally, such a category that

moreover has a dagger (see [2]) could hopefully embed into a category of inner product semimodules.

Acknowledgement

This work evolved from the elementary observation $\mathbf{cSRng} = \mathbf{cMon}(\mathbf{cMon})$ that is due to Bart Jacobs, who moreover made helpful suggestions.

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A The Hilton-Eckmann argument

The following argument is due to Hilton and Eckmann ([4], but see also [15, Exercise II.5.5]), who proved it for $\mathbf{C} = \mathbf{Set}$. It states that when an object carries two monoid structures and the multiplication map of one is a monoid homomorphism with respect to the other, then the two monoid structures coincide and are in fact commutative.

Theorem A.1 (Hilton-Eckmann) *Let (\mathbf{C}, \otimes, I) be a symmetric monoidal category. Suppose given an object $X \in \mathbf{C}$ and morphisms $\mu_1, \mu_2 : X \otimes X \rightarrow X$ and $\eta_1, \eta_2 : I \rightarrow X$. If (X, μ_1, η_1) and (X, μ_2, η_2) are both monoids (in \mathbf{C}) and the following diagram commutes,*

$$\begin{array}{ccc}
 X \otimes X \otimes X \otimes X & \xrightarrow{\mu_2 \otimes \mu_2} & X \otimes X \\
 \text{id} \otimes \gamma \otimes \text{id} \downarrow \cong & & \downarrow \mu_1 \\
 X \otimes X \otimes X \otimes X & & X \\
 \mu_1 \otimes \mu_1 \downarrow & & \\
 X \otimes X & \xrightarrow{\mu_2} & X
 \end{array} \tag{A.1}$$

then $(X, \mu_1, \eta_1) = (X, \mu_2, \eta_2)$ is in fact a commutative monoid (in \mathbf{C}).

Proof. First we show that $\eta_1 = \eta_2$.

To prevent a forest of diagrams, we give rest of the proof for $\mathbf{C} = \mathbf{Set}$. The reader can check for herself that it generalizes to any symmetric monoidal category. Let us further abbreviate $\eta_1 = \eta_2$ to 1, $\mu_1(x, y)$ to $x \circ y$, and $\mu_2(x, y)$ to $x \bullet y$.

$$\begin{aligned}
 x \circ y &= (1 \bullet x) \circ (y \bullet 1) \stackrel{(A.1)}{=} (1 \circ y) \bullet (x \circ 1) = y \bullet x \\
 &= (y \circ 1) \bullet (1 \circ x) \stackrel{(A.1)}{=} (y \bullet 1) \circ (1 \bullet x) = y \circ x.
 \end{aligned}$$

□